

COUNTABLE CATEGORICITY

2. OLIGOMORPHIC PERMUTATION GROUPS

2.1. Topology. Let \mathbb{B} be a relational structure. We define $\text{Sym}(B)$ to be the set of all permutations of the domain B of \mathbb{B} .

Definition 1. We say that $\mathcal{P} \subseteq \text{Sym}(\mathbb{B})$ is *closed* if the following condition holds:

For all $\alpha \in \text{Sym}(B)$ if for every finite $A \subseteq B$ there is some $\beta \in \mathcal{P}$ such that for all $x \in A$ we have that $\alpha(x) = \beta(x)$ then $\alpha \in \mathcal{P}$.

Remark 2. Let us say that $\mathcal{Q} \subseteq \text{Sym}(B)$ is *open* if its complement is closed. Then, by unfolding the definitions, we get that \mathcal{Q} is open if, and only if, it is a union of sets of the form:

$$\{\alpha \in \text{Sym}(B) : (\forall x \in A)\alpha(x) = \beta(x)\},$$

for some fixed $\beta \in \text{Sym}(B)$ and finite set $A \subseteq B$ (a set of this form is the *coset of the stabilisers of a finite tuple*).

Definition 3. Let $\mathcal{P} \subseteq \text{Sym}(B)$. We define the *strong invariants* of \mathcal{P} , denoted $\text{slnv}(\mathcal{P})$ to be the set of all relations R on B such that for all $\alpha \in \mathcal{P}$ both α and α^{-1} preserve R .

Proposition 4. Let $\mathcal{P} \subseteq \text{Sym}(B)$. Then, the following are equivalent:

- (1) \mathcal{P} is the automorphism group of some relational structure with domain B .
- (2) \mathcal{P} is a closed subgroup of $\text{Sym}(B)$.
- (3) \mathcal{P} is the automorphism group of a homogeneous relational structure with domain B

Proof.

- (1) \Rightarrow (2): Of course, if \mathcal{P} is an automorphism group it is a group, so we need only show that it is closed. Suppose that $\alpha \in \text{Sym}(B)$ is such that for every finite $A \subseteq B$ there is some $\beta \in \mathcal{P}$ such that $\alpha x = \beta x$ for all $x \in A$. Then, α must be itself an automorphism, for otherwise this would be witnessed from the restriction of α to a finite set.
- (2) \Rightarrow (3): We claim that if \mathcal{P} is a closed subgroup of $\text{Sym}(B)$ then $\text{slnv}(\mathcal{P})$ is homogeneous and $\mathcal{P} = \text{Aut}(\text{slnv}(\mathcal{P}))$. First, we obviously have that $\mathcal{P} \subseteq \text{Aut}(\text{slnv}(\mathcal{P}))$. Now, suppose that $\alpha \in \text{Aut}(\text{slnv}(\mathcal{P}))$. Then, for all finite $\{a_1, \dots, a_n\} \subseteq B$ consider the relation:

$$R_{a_1, \dots, a_n} = \{(\beta a_1, \dots, \beta a_n) : \beta \in \mathcal{P}\}.$$

Clearly, R is in $\text{slnv}(\mathcal{P})$, since \mathcal{P} is a group, and of course, since α is an automorphism of $\text{slnv}(\mathcal{P})$ then it must preserve this set. In particular, there is some $\beta \in \mathcal{P}$ such that $\alpha a_i = \beta a_i$ for all $i = 1, \dots, n$. Since \mathcal{P} is closed, we must have that $\alpha \in \mathcal{P}$. We actually have also essentially showed homogeneity, since, any finite isomorphism, say with domain $\{a_1, \dots, a_n\}$ preserves the relation R_{a_1, \dots, a_n} , and thus (again by closure) this finite isomorphism is the restriction of an element in \mathcal{P} .

- (3) \Rightarrow (1): Trivial.

□

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Remark 5. We have seen that being ω -categorical is a property of the automorphism group. Thus, in fact by the previous proposition we have that every ω -categorical structure \mathbb{B} has the same automorphism group as a homogeneous ω -categorical structure, which, by the proof, is just $\text{slnv}(\text{Aut}(\mathbb{B}))$.

2.2. The slnv -Aut Galois connection.

Definition 6. An (*anti-tone*) Galois connection is a pair of functions:

$$F : U \rightarrow V, \text{ and } G : V \rightarrow U,$$

between two posets F and G such that:

$$v \leq F(u) \text{ if, and only if, } u \leq G(v),$$

for all $u \in U$ and $v \in V$.

Remark 7. For all $u \in U$ and $v \in V$ we have the following:

- $u \leq G(F(u))$ [Why? Apply the definition to $F(u) \leq F(u)$].
- $v \leq F(G(v))$
- $F(u) = F(G(F(u)))$ [Why? We have already shown \leq , for \geq apply the definition to $u \leq G(F(u)) \leq G(F(G(F(u))))$]
- $G(v) = G(F(G(v)))$

Proposition 8. Let B be a set. The operations Aut and slnv form a Galois connection between the set of all relations over B and the sets of permutations of B .

Proof. Let \mathcal{R} be a set of relations over B and \mathcal{P} a set of permutations of B . We have to show that:

$$\mathcal{P} \subseteq \text{Aut}(\mathcal{R}) \text{ if, and only if, } \mathcal{R} \subseteq \text{slnv}(\mathcal{P}).$$

\Rightarrow : If $\mathcal{P} \subseteq \text{Aut}(\mathcal{R})$ then for all $R \in \mathcal{R}$ and $g \in \mathcal{P}$ we have that both g and g^{-1} preserve R .

\Leftarrow : If $\mathcal{R} \subseteq \text{slnv}(\mathcal{P})$ then for all $g \in \mathcal{P}$ both g and g^{-1} preserve R , so $g \in \text{Aut}(\mathcal{R})$. □

Definition 9. Let $\mathcal{P} \subseteq \text{Sym}(B)$. We define the following:

- The *permutation group generated by \mathcal{P}* , denoted $\langle \mathcal{P} \rangle$, to be the smallest permutation group containing \mathcal{P} .
- The *closure of \mathcal{P} in $\text{Sym}(B)$* , denoted $\overline{\mathcal{P}}$, to be the smallest closed subset of $\text{Sym}(B)$ containing \mathcal{P} .

Remark 10. Explicitly, $\overline{\mathcal{P}}$ contains \mathcal{P} together with all the permutations $\alpha \in \text{Sym}(B)$ such that for all finite $A \subseteq B$ there is some $\beta \in \mathcal{P}$ such that $\alpha(x) = \beta(x)$ for all $x \in A$ (i.e. it is exactly \mathcal{P} together with its “limit points”).

Proposition 11. Let $\mathcal{P} \subseteq \text{Sym}(B)$, and define \mathcal{P}^* to be the smallest permutation group that contains \mathcal{P} and is closed in $\text{Sym}(B)$. Then:

$$\mathcal{P}^* = \overline{\langle \mathcal{P} \rangle} = \text{Aut}(\text{slnv}(\mathcal{P})).$$

Proof.

- First we shall show that $\mathcal{P}^* = \overline{\langle \mathcal{P} \rangle}$. Note that $\mathcal{P} \subseteq \mathcal{P}^*$ and since \mathcal{P}^* is a permutation group we immediately get (by minimality) that $\langle \mathcal{P} \rangle \subseteq \mathcal{P}^*$. Now, since \mathcal{P}^* is closed, this (again by minimality) implies that $\overline{\langle \mathcal{P} \rangle} \subseteq \mathcal{P}^*$. Let us then show the converse. Since \mathcal{P}^* is the smallest permutation group that contains \mathcal{P} and is closed in $\text{Sym}(B)$ (and $\overline{\langle \mathcal{P} \rangle}$ contains \mathcal{P}) it suffices to show that $\overline{\langle \mathcal{P} \rangle}$ is a permutation group (we already know

that it is closed). Showing that it contains the identity is immediate (since $\langle \mathcal{P} \rangle$ is a permutation group), so it suffices to show closure under inverses and composition. Let us only do the former (the argument for the latter is similar). Suppose that $\alpha, \beta \in \overline{\langle \mathcal{P} \rangle}$. By definition of $\overline{\langle \mathcal{P} \rangle}$ we must have that for all finite $A \subseteq \beta$ there are $\alpha', \beta' \in \langle \mathcal{P} \rangle$ such that $\alpha x = \alpha' x$ and $\beta x = \beta' x$ for all $x \in A$. Since $\langle \mathcal{P} \rangle$ is a permutation group it contains their composition, and hence, since $\overline{\langle \mathcal{P} \rangle}$ is closed it must contain $\alpha\beta$.

- Now we shall show that $\overline{\langle \mathcal{P} \rangle} = \text{Aut}(\text{slnv}(\mathcal{P}))$. On the one hand, if $\alpha \in \overline{\langle \mathcal{P} \rangle}$, then we claim that both α and α^{-1} preserve each $R \in \text{slnv}(\mathcal{P})$. Indeed, suppose that $R \in \text{slnv}(\mathcal{P})$ and $t \in R$. By definition of $\overline{\langle \mathcal{P} \rangle}$ there are $\beta_1, \dots, \beta_k \in \mathcal{P} \cup \mathcal{P}^{-1}$ such that $\alpha t = (\beta_1 \circ \dots \circ \beta_k)t$. Since each β_i preserves R it follows that so does α (and α^{-1} , similarly). Conversely, if $\alpha \in \text{Aut}(\text{slnv}(\mathcal{P}))$ then α and its inverse preserve the relation:

$$\{(\beta t_1, \dots, \beta t_n) : \beta \in \langle \mathcal{P} \rangle\},$$

and by closure we must have that $\alpha \in \overline{\langle \mathcal{P} \rangle}$. □

Proposition 12. *Let \mathbb{B} be any structure. Then:*

$$\langle \mathbb{B} \rangle_{\text{fo}} \subseteq \text{slnv}(\text{Aut}(\mathbb{B})),$$

where $\langle \mathbb{B} \rangle_{\text{fo}}$ denotes the set of all first-order definable relations in \mathbb{B} .

Proof. If $R \in \langle \mathbb{B} \rangle_{\text{fo}}$ then all $g, g^{-1} \in \text{Aut}(\mathbb{B})$ preserve R . □

The following is an immediate consequence of Ryll-Nardzewski:

Theorem 13. *If \mathbb{B} is ω -categorical, in a countable signature, then:*

$$\text{slnv}(\text{Aut}(\mathbb{B})) = \langle \mathbb{B} \rangle_{\text{fo}},$$

and this characterises ω -categorical structures.

Remark 14. So, $\text{Aut}(\mathbb{B})$ is precisely the automorphism group of $\text{slnv}(\text{Aut}(\mathbb{B}))$, which in the ω -categorical case is precisely $\langle \mathbb{B} \rangle_{\text{fo}}$. Thus, if \mathbb{B} is ω -categorical we have that:

$$\text{Aut}(\mathbb{B}) = \text{Aut}(\langle \mathbb{B} \rangle_{\text{fo}}).$$

Theorem 15. *Let \mathbb{B} be a countable ω -categorical structure. Then:*

- *The sets of the form $\langle \mathbb{A} \rangle_{\text{fo}}$ where \mathbb{A} is a first-order reduct of \mathbb{B} ordered by inclusion form a lattice.*
- *The closed supergroups of $\text{Aut}(\mathbb{B})$ in $\text{Sym}(B)$ ordered by inclusion form a lattice.*
- *The operator slnv is an anti-isomorphism between these two lattices and Aut is its inverse.*

REFERENCES

- [Bod21] Manuel Bodirsky. *Complexity of Infinite-Domain Constraint Satisfaction*. Cambridge University Press, May 2021.