

CSP Reading Group Sections 6.1.1–6.1.3

B set, $n \geq 1$ $O_B^{(n)} = B^{B^n} = \{ \text{maps } B^n \rightarrow B \}$... n-ary operations on B

$$O_B = \bigcup_{n \geq 1} O_B^{(n)} \quad \dots \text{operations on } B$$

operation clone : $\mathcal{Q} \subseteq O_B$ satisfying

- ↓
sometimes called "concrete clone"
- \mathcal{Q} contains all projections $\Pi_i^n \forall n \in \{1, \dots, n\}$
 - \mathcal{Q} closed under composition, i.e.
 $\forall f \in \mathcal{Q}$ n-ary and $\forall g_1, \dots, g_n \in \mathcal{Q}$ k-ary
 $f(g_1, \dots, g_n) \in \mathcal{Q}$ *→ this gives in fact every composition*
 $(x_1, \dots, x_k) \mapsto f(g_1(x_1, \dots, x_k), \dots, g_n(x_1, \dots, x_k))$

$y \subseteq O_B \rightsquigarrow \langle y \rangle \dots$ smallest clone that contains y

DEF.:

polymorphism of a structure \mathbb{B} is a homomorphism from a finite power of \mathbb{B} to \mathbb{B}

$\text{Pol}(\mathbb{B}) = \text{the clone of all polymorphisms of } \mathbb{B}$
easy to verify

$f \in O_B^{(n)}$ preserves $R \subseteq B^k$ if $\forall r_1, \dots, r_n \in R$

\uparrow
(R is invariant under f)

$$f((\underset{r_1 \in R}{\underset{\dots}{\dots}}, \underset{r_n \in R}{\underset{\dots}{\dots}})) = \begin{pmatrix} f(r_1) \\ \vdots \\ f(r_n) \end{pmatrix} \in R$$

OBSERVATION: f is a polymorphism of $\mathbb{B} \Leftrightarrow f$ preserves all relations of \mathbb{B}

DEF.: operation clone $\mathcal{L} \subseteq \mathcal{O}_B$ is finitely related if
 \exists structure B with finite relational signature
s.t. $\text{Pol}(B) = \mathcal{L}$

$$R_B = \bigcup_{n \geq 1} R_B^{(n)} \dots \text{ (n-ary) relations on } B$$

NOTATION: $R \subseteq R_B \rightsquigarrow \text{Pol}(R)$
 $\gamma \subseteq \mathcal{O}_B \rightsquigarrow \text{Inv}(\gamma)$

PROPOSITION: For every B $\underbrace{\langle B \rangle_{\text{pp}}}_{\text{relations definable by a pp-formula}} \subseteq \text{Inv}(\text{Pol}(B))$

Proof: $R \subseteq B^k$ $f \in \text{Pol}(B)$ l -ary

pp-def. of R $\Psi(x_1, \dots, x_k) = \exists x_{k+1}, \dots, x_n \quad \Psi'(x_1, \dots, x_{k+1}, x_{k+1}, \dots, x_n)$
 $t_1, \dots, t_l \in R \rightsquigarrow s_1, \dots, s_l$ extended by witnesses so that they satisfy
 $\Psi' \Rightarrow f(s_1, \dots, s_l)$ satisfies $\Psi' \Rightarrow f(t_1, \dots, t_l) \in R$ ■

6.1.1 Pol-Inv

$B \dots$ countably infinite

DEF : $\mathcal{Y} \subseteq \mathcal{O}_B$ is locally closed \Leftrightarrow

$[\forall n \geq 1 \ \forall g \in \mathcal{O}_B^{(n)} : (\exists F \subseteq B \text{ finite} \ \exists f \in \mathcal{Y} \ g|_F = f|_F) \Rightarrow g \in \mathcal{Y}]$

$\overline{\mathcal{Y}}$... local closure of \mathcal{Y} = smallest locally closed subset containing \mathcal{Y}

$f \in \overline{\langle \mathcal{Y} \rangle}$... f is locally generated by \mathcal{Y}

PROPOSITION : $\mathcal{C} \subseteq \mathcal{O}_B$... the smallest locally closed clone

that contains $\mathcal{Y} \subseteq \mathcal{O}_B \Rightarrow \underline{\mathcal{C} = \overline{\langle \mathcal{Y} \rangle} = \text{Pol}(\text{Inv}(\mathcal{Y}))}$.

PROPOSITION: $\mathcal{C} \subseteq \mathcal{O}_B$... the smallest locally closed clone

that contains $\mathcal{G} \subseteq \mathcal{O}_B \Rightarrow \mathcal{C} = \overline{\langle \mathcal{G} \rangle} = \text{Pol}(\text{Inv}(\mathcal{G}))$.

Proof:

$\overline{\langle \mathcal{G} \rangle} \subseteq \mathcal{C}$ is clear because \mathcal{C} is a locally closed clone

For $\mathcal{C} \subseteq \overline{\langle \mathcal{G} \rangle}$ we show that $\overline{\langle \mathcal{G} \rangle}$ is a locally closed clone
(since $\mathcal{G} \subseteq \overline{\langle \mathcal{G} \rangle}$).

$\overline{\langle \mathcal{G} \rangle}$ is locally closed and contains projections.

$\overline{\langle \mathcal{G} \rangle}$ is closed under composition: $f_1 g_1, \dots, g_n \in \overline{\langle \mathcal{G} \rangle}$, $F \subseteq B$ finite
approximate g_1, \dots, g_n on F by $g'_1, \dots, g'_n \in \langle \mathcal{G} \rangle$ and
 f by f' on $g_1(F) \cup \dots \cup g_n(F)$, then $f'(g'_1, \dots, g'_n)|_F = f(g_1, \dots, g_n)|_F$
 $\Rightarrow \overline{\langle \mathcal{G} \rangle} = \mathcal{C}$

$\langle \Psi \rangle \subseteq \text{Pol}(\text{Inv}(\Psi))$:

$f \in \langle \Psi \rangle$ n-ary, $t_1, \dots, t_n \in R$

Find $g \in \langle \Psi \rangle$ s.t. $f(t_1, \dots, t_n) = g(t_1, \dots, t_n)$.

$\text{Inv}(\Psi) = \text{Inv}(\langle \Psi \rangle) \Rightarrow g \in \langle \Psi \rangle \subseteq \text{Pol}(\text{Inv}(\langle \Psi \rangle)) = \text{Pol}(\text{Inv}(\Psi))$

$\Rightarrow f(t_1, \dots, t_n) = g(t_1, \dots, t_n) \in R$

$\text{Pol}(\text{Inv}(\Psi)) \subseteq \langle \Psi \rangle$:

$f \in \text{Pol}(\text{Inv}(\Psi))$ n-ary, $F \subseteq B$ finite.

$t_1, \dots, t_m \dots$ all elements of F^n

$R := \{ (g(t_1), \dots, g(t_m)) \mid g \in \langle \Psi \rangle^{(n)} \}$ is in $\text{Inv}(\Psi)$

$\Rightarrow f$ preserves R

$\prod_{i=1}^m \in \langle \Psi \rangle^{(n)}$ $\forall i \Rightarrow (t_1[i], \dots, t_m[i]) \in R \quad \forall i \Rightarrow f(t_1, \dots, t_m) \in R$

$\Rightarrow (f(t_1), \dots, f(t_m)) = (g(t_1), \dots, g(t_m))$ for some $g \in \langle \Psi \rangle \Rightarrow f|_F = g|_F$ ■

COROLLARY: For $\mathcal{Y} \subseteq \mathcal{O}_B$ TFAE:

- (1) \mathcal{Y} is the polymorphism clone of a relational structure
- (2) \mathcal{Y} is a locally closed clone

Proof:

(1) \Rightarrow (2): $\mathcal{Y} = \text{Pol}(\mathbb{A})$... we have seen that these are locally closed and clones

(2) \Rightarrow (1): $\mathcal{Y} = \overline{\langle \mathcal{Y} \rangle} = \underbrace{\text{Pol}(\text{Inv}(\mathcal{Y}))}_{\text{gives rise to a structure on } B}$ ■

Locally closed subclones of \mathcal{O}_B form a complete lattice:

$$(\mathcal{Y}_i)_{i \in I} \begin{array}{l} \xrightarrow{\text{join}} \overline{\left\langle \bigcup_{i \in I} \mathcal{Y}_i \right\rangle} \\ \xrightarrow{\text{meet}} \bigcap_{i \in I} \mathcal{Y}_i \end{array}$$

DEF: G ... a permutation group on the set B

We write $\text{Orb}(G)$ for a relational structure on B
whose relations are orbits of n -tuples of G $\forall n \geq 1$.

Some properties:

- $\text{Pol}(\text{Orb}(G)) =$ the unique largest clone ℓ s.t. $\ell^{(1)} = \overline{G}$
- $\text{End}(\text{Orb}(G)) = \overline{G}$
- G oligomorphic \Rightarrow (fo-expansions of) $\text{Orb}(G)$ are mc-cores

typo in the book
(I think)

REMARK: $\text{Pol}(\text{Orb}(G))^{(n)}$ consists precisely of operations
 $f \in \mathcal{O}_B^{(n)}$ s.t. $\forall g_1, \dots, g_n \in G$
 $x \mapsto f(g_1(x), \dots, g_n(x))$ is in \overline{G}

6.1.2 Inv-Pol

DEF.: infinitary pp-formulas are defined inductively:

- atomic formulas
- finite and infinite conjunctions of formulas $\phi(x_1, \dots, x_n)$ in the class with the same free variables x_1, \dots, x_n
- $\exists x_i \Phi(x_1, \dots, x_n)$ where $\Phi(x_1, \dots, x_n)$ is in the class

EXAMPLE: $\mathbb{B} = (\mathbb{Z} \cup \{\infty\}, <)$

$\forall i \in \mathbb{N} \quad \{(n, m) \in \mathbb{B}^2 \mid n + i < m\}$ is pp-definable by $\phi_i(x, y)$.
 $\{\infty\}$ by $\exists x \bigwedge_{i=1}^{\infty} \phi_i(x, y)$.

pp-formulas satisfiable by ∞ are always satisfiable by $z \in \mathbb{Z} \Rightarrow$ all elements of \mathbb{Z} \Rightarrow infinite conjunctions of pp-formulas always define a relation containing \mathbb{Z}

chain: $(R_i)_{i \in \mathbb{N}}$ all R_i of the same arity and $R_i \subseteq R_{i+1} \cup t_i$

THEOREM: \mathbb{B} countable

$R \subseteq B^n$ is preserved by $\text{Pol}(\mathbb{B}) \iff R$ is the union of a chain of relations that have infinitary pp-definitions in \mathbb{B}

Proof:

Easy to check that every such R is preserved by $\text{Pol}(\mathbb{B})$
(we have already seen for classical pp-definitions).

$\text{RelInv}(\text{Pol}(\mathbb{B}))$ of arity n .

Let a_1, a_2, \dots be an enumeration of R .

$R_i :=$ intersection of all relations with infinitary pp-definitions
in \mathbb{B} containing $a_{1, \dots, i}$

$R_i \subseteq R_{i+1} \cup t_i$

We show $R = \bigcup_{i \in \mathbb{N}} R_i$.

$R \subseteq \bigcup_{i \in \mathbb{N}} R_i$: Clear (B^n is always in the intersections).

To prove \supseteq , we show $R_i \subseteq R$ f.t., let $t = (t_1, \dots, t_n) \in R_i$.

We search for $f \in \text{Pol}(B)$ of arity i s.t. $f(a_1, \dots, a_i) = \binom{t_1}{t_n}$.

b_1, b_2, \dots enumeration of B^i s.t.

WLOG no rows repeat, otherwise
we find a definition of a projection
of R and then pp-define R
from it by equalities

$$\begin{array}{c} b_1 \rightarrow \\ b_2 \rightarrow \\ \vdots \\ b_n \rightarrow \end{array} \left(\begin{array}{c|c|c|c} a_1 & a_2 & \dots & a_i \\ \hline \vdots & & & \vdots \end{array} \right)$$

Define $f(b_1) = t_1, \dots, f(b_n) = t_n \rightsquigarrow$ this partial definition preserves
all infinitary pp-formulas since
 $t \in R_i$

Extend f already defined on $b_1, \dots, b_m \in B^i$, $m \geq n$:
 Take $\phi(\bar{x}_1, \dots, \bar{x}_{m+1})$ conjunction of all infinitary pp-formulas satisfied by b_1, \dots, b_{m+1}

f preserves $\phi'(\bar{x}_1, \dots, \bar{x}_m) = \exists \bar{x}_{m+1}: \phi(\bar{x}_1, \dots, \bar{x}_{m+1})$
 $\Rightarrow B \models \phi'(f(b_1), \dots, f(b_m)) \rightsquigarrow t_{m+1}$ witness for ϕ'
 $f(b_{m+1}) := t_{m+1}$

this extension preserves every infinitary pp-formula over B
 in particular, $f \in \text{Pol}(B) \Rightarrow$

$$t = \left(\begin{array}{c} f(b_1) \\ \vdots \\ f(b_n) \end{array} \right) = f\left(\underbrace{a_1}_{\in R} \mid \dots \mid \underbrace{a_i}_{\in R} \right) \in R$$

$$\Rightarrow R_i \subseteq R \quad \stackrel{i: \text{arbitrary}}{\Rightarrow} \quad R = \bigcup_{i \in \mathbb{N}} R_i \quad \text{as we wanted} \quad \blacksquare$$

6.1.3 Oligomorphic clones

$\mathcal{C} \subseteq \mathcal{O}_B$ operation clone

$e \in \mathcal{C}^{(1)}$ is invertible if $\exists i \in \mathcal{C}^{(1)} : \forall x \in B \quad e(i(x)) = i(e(x)) = x$

$\mathcal{C} = \text{Pol}(B) \Rightarrow$ invertible operations in \mathcal{C} are $\text{Aut}(B)$

DEF.: $\mathcal{C} \subseteq \mathcal{O}_B$... operational clone

\mathcal{C} is oligomorphic if the set of invertible operations in \mathcal{C} forms an oligomorphic permutation group.

THM 4.1.6 + COR 6.1.6 imply:

locally closed clone $\mathcal{C} \subseteq \mathcal{O}_B$ is oligomorphic



$\mathcal{C} = \text{Pol}(B)$, B ω -categorical

THEOREM: \mathbb{B} countable w-categorical
 $\underbrace{\langle \mathbb{B} \rangle_{\text{pp}}}_{\text{relations}} = \text{Inv}(\text{Pol}(\mathbb{B}))$
pp-definable in \mathbb{B}

Proof: \subseteq we already know

\supseteq : For each fixed arity ... only finitely many pp-definable relations

\Rightarrow infinite conjunctions of pp-formulas are equivalent
to pp-formulas

\Rightarrow unions of chains of pp-definable relations are pp-def.

THEOREM
 \Rightarrow $R \in \text{Inv}(\text{Pol}(\mathbb{B}))$ is pp-definable



THEOREM: B countable, w-categorical

- (1) sets $\langle \mathcal{C} \rangle_{pp}$, \mathcal{C} fo-reducts of B , form a lattice
- (2) closed subclones of O_B containing $\text{Aut}(B)$ form a lattice
- (3) operators Inv and Pol are mutually inverse anti-isomorphisms of these lattices

COROLLARY: B w-categorical, $\sigma f \in O_B$

- $\langle \sigma f \rangle = \text{Pol}(B)$ iff $\langle B \rangle_{pp} = \text{Inv}(\sigma f)$
- the smallest relation in $\langle B \rangle_{pp}$ containing $R \subseteq B^k$ is $\{f(a_1, \dots, a_n) \mid n \in \mathbb{N}, f \in \text{Pol}^{(n)}(B), a_1, \dots, a_n \in R\}$

EXAMPLE: $\mathbb{B} = (\mathbb{Z}, \{(x,y) \mid x = y + 1\})$

$\text{Inv}(\text{Pol}(\mathbb{B})) = \langle \mathbb{B} \rangle_{\text{pp}}$ (even though \mathbb{B} is not ω -cat.)

We show $\langle \mathbb{B} \rangle_{\text{pp}}$ is closed under infinite conjunctions and unions of chains.

Fact: Every relation is a conjunction of binary relations.

Binary relations: $\emptyset, \mathbb{Z}^2, \{(x,y) \mid x = y + c\}, c \in \mathbb{Z}$

\hookrightarrow infinite conjunctions of these are of the same form

\hookrightarrow unions of such chains are of this form